



Principles and Practice of Application Performance Measurement and Analysis on Parallel Systems

Lecture 1: Terminology and Methodology

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Performance Tuning: an Old Problem!



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"The most constant difficulty in contriving the engine has arisen from the desire to reduce the time in which the calculations were executed to the shortest which is possible."

> Charles Babbage 1791 - 1871

Motivation



- High complexity in parallel and distributed systems
 - □ Four layers
 - Application
 - □ Algorithm, data structures
 - Parallel programming interface / Middle ware
 - □ Compiler, parallel libraries, communication, synchronization
 - Operating system
 - □ Process and memory management, IO
 - Hardware
 - □ CPU, memory, network
- Mapping/interaction between different layers

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Performance Properties of Parallel Programs

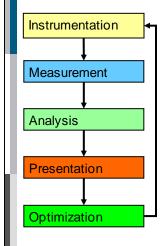


- Factors which influence performance of parallel programs
 - ☐ "Sequential" factors
 - Computation
 - ⇒ Choose right algorithm, use optimizing compiler
 - Cache and memory
 - ⇒ Tough! Not many tools yet, hope compiler gets it right
 - Input / output
 - ⇒ Not given enough attention
 - □ "Parallel" factors
 - Communication (Message passing)
 - Threading
 - Synchronization
 - ⇒ More or less understood, tool support

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Performance Measurement Cycle





- Insertion of extra code (probes, hooks) into application
- Collection of data relevant to performance analysis
- Calculation of metrics, identification of performance problems
- Transformation of the results into a representation that can be easily understood by a human user
- Elimination of performance problems



CONTENT Metrics Instrumentation techniques Source code instrumentation Binary instrumentation Instrumentation of parallel programs MPI OpenMP Measurement techniques Profiling Tracing

Metrics of Performance What can be measured? A count of how many times an event occurs E.g., Number of input / output requests The duration of some time interval E.g., duration of these requests The size of some parameter Number of bytes transmitted or stored Derived metrics E.g., rates / throughput Needed for normalization



JÜLICH Example Metrics Clock rate MIPS ☐ Millions of instructions executed per second ☐ Floating-point operations per second Benchmarks ☐ Standard test program(s) 'math" Operations? **HW Operations?** ☐ Standardized methodology HW Instructions? □ E.g., SPEC, Linpack 32-/64-bit?.. QUIPS / HINT [Gustafson and Snell, 95] □ Quality improvements per second □ Quality of solution instead of effort to reach it ■ Execution time 1-9

Execution Time

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- Wall-clock time
 - $\hfill\square$ Includes waiting time: IO, memory, other system activities
 - □ In time-sharing environments also time consumed by other applications
- CPU time
 - ☐ Time spent by the CPU to execute the program
 - □ Execution time on behalf of the program
 - $\hfill\square$ Does not include time the program was context-switched out
 - Problem: does not include inherent waiting time (e.g., IO)
 - Problem: portability? What is user, what is system time?
- Problem: execution time is non-deterministic
 - ☐ Use mean or minimum of several runs



Load Imbalance Metrics



- Imbalance Time
 - ☐ Metric time to identify code regions that need optimization
 - □ Two variations:
 - ☐ Computation Imbalance Time

Computation Imbalance Time = Max Time - Avg time

☐ Synchronization Imbalance Time

Synchronization Imbalance Time = Avg Time - Min time

- □ Provides an estimation to the user of how much time in the overall program would be saved if the corresponding section of the code had a perfect balance
 - Represents an upper bound on the "potential saving"

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Load Imbalance Metrics



- Imbalance %
 - ☐ Provide an idea of the "badness" of the imbalance
 - ☐ Corresponds to the % of the time that the rest of the team, excluding the slowest PE is not engaged in useful work on the given function
 - "Percentage of resources available for parallelism" that is wasted

Imbalance% =
$$100 \text{ X}$$
 Imbalance time $\frac{N}{N-1}$



JÜLICH Speedup and Efficiency ■ For a given problem A, let □ SerTime(n) = Time of the best serial algorithm to solve A for input of size n = Time of the parallel algorithm + architecture □ ParTime(n,p) to solve A for input size n, using p processors \Box Note that SerTime(n) ≤ ParTime(n,1) Then □ Speedup(p) = SerTime(n) / ParTime(n,p) □ Work(p) = p • ParTime(n,p) □ Efficiency(p) = SerTime(n) / [p • ParTime(n,p)] I-13

Speedup and Efficiency II



■ In general,	expect		
		0	≤ Spe

 $eedup(p) \le p$

Serial work ≤ Parallel work < ∞

≤ 1 0 ≤ Efficiency

■ Linear speedup: if there is a constant c > 0 so that speedup is at least $c \cdot p$. Many use this term to mean c = 1.

■ Perfect or ideal speedup: speedup(p) = p

■ Superlinear speedup: speedup(p) > p (effiency > 1)

☐ Typical reason: Parallel computer has p times more memory (cache), so higher fraction of program data fits in memory instead of disk (cache instead of memory)

☐ Parallel version is solving slightly different, easier problem or provides slightly different answer



Amdahl's Law

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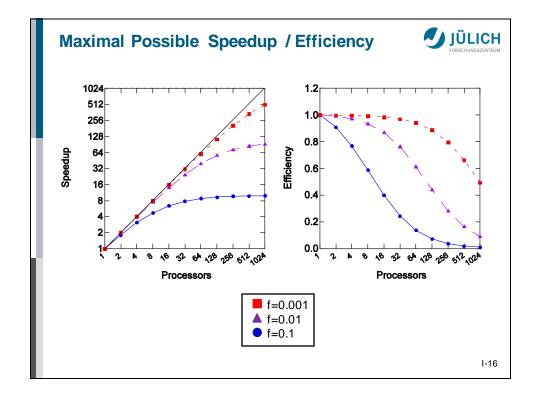
- Amdahl [1967] noted:
 - ☐ Given a program, let f be fraction of time spent on operations that must be performed serially (unparallelizable work). Then for p processors:

Speedup(p)
$$\leq \frac{1}{f + (1 - f)/p}$$

☐ Thus no matter how many processors are used

Speedup(p) $\leq 1/f$

☐ Unfortunately, typical f is 5 – 20%





Amdahl's Law II



- Amdahl was an optimist
 - □ Parallelization might require extra work, typically
 - Communication
 - Synchronization
 - Load balancing
 - ☐ Amdahl convinced many people that general-purpose parallel computing was not viable
- Amdahl was an pessimist
 - □ Fortunately, we can break the law!
 - □ Find better (parallel) algorithms with much smaller values of f
 - □ Superlinear speedup because of more data fits cache/memory
 - □ Scaling: time spent in serial portion is often a decreasing fraction of the total time as problem size increase

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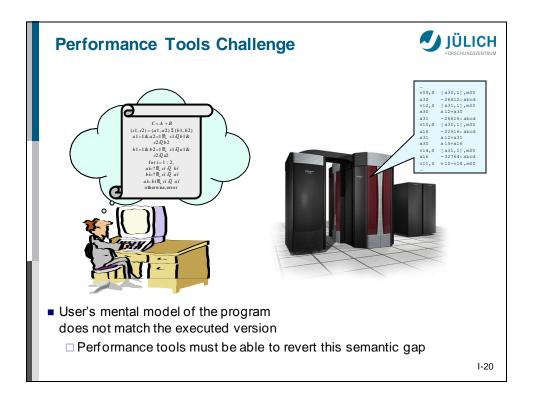
Scaling



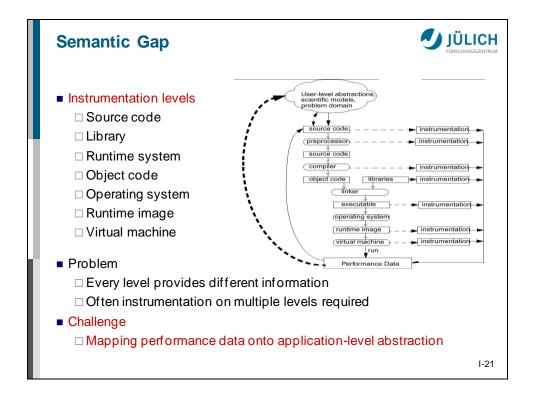
- Sometimes the serial portion
 - ☐ is a fixed amount of time independent of problem size
 - □ or grows with problem size but slower than total time
- Thus can often exploit large parallel machines by scaling the problem size with the number of processes
- Scaling approaches used for speedup reporting/measurements:
 - ☐ Fixed problem size (⇒ strong scaling)
 - ☐ Fixed problem size per processor (☐ weak scaling)
 - ☐ Fixed time, find largest problem solvable [Gustafson 1988] Commonly used in evaluating databases (transactions/s)
 - □ Fixed efficiency: find smallest problem to achieve it
 (⇒ isoefficiency analysis)











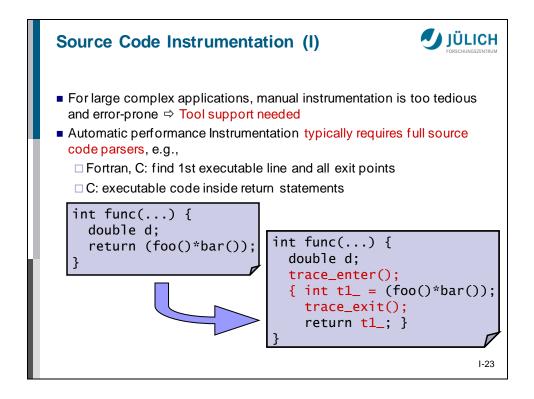
Instrumentation Techniques



- Static instrumentation
 - ☐ Program is instrumented prior to execution
- Dynamic instrumentation
 - □ Program is instrumented at runtime
- Code is inserted
 - Manually
 - □ Automatically
 - By preprocessor / source-to-source translation tool
 - By compiler
 - By linking against pre-instrumented library or runtime system
 - By binary-rewrite / dynamic instrumentation tool
- ⇒ e.g., "printf" ⇒ manual static source-code instrumentation

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JÜLICH Source Code Instrumentation (II) ■ Example C++ issues: □ Template instrumentation? □ Function overloading □ Executing code before main □ Operator overloading ■ C++ instrumentation trick □ Define instrumentation object class Tracer { public: Tracer(...) { trace_enter(); } ~Tracer() { trace_exit(); } □ Declare instrumentation object as 1st statement in every function and method to be instrumented int func(...) { Tracer trc_1; double d; return (foo()*bar()); 1-24



TAU Source Code Instrumentor



- Part of the TAU performance framework
- Supports
 - □ f77, f90
 - □ C, and C++
- Inserts calls to the TAU monitoring API
- Based on the Program Database Toolkit





http://tau.uoregon.edu/

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Program Database Toolkit



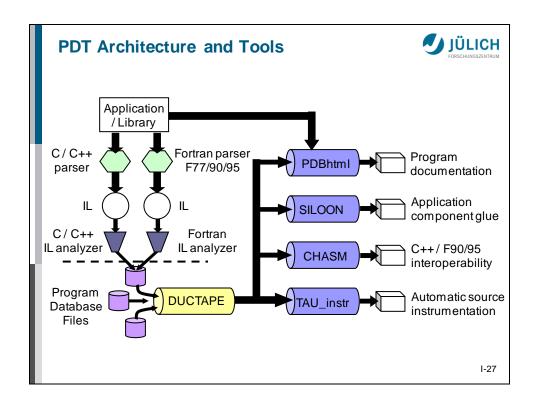
- Based on commercial parsers
- □ C, C++: Edison Design Group (EDG)
 - Full ISO 1998 C++ and ISO 1999 C Support
 - □ Fortran 77, Fortran90: Mutek, [Cleanscape]



Program Database Utilities and Conversion Tools APplication Environment (DUCTAPE)

- ☐ Object-oriented Access to Static Information
- □ Classes, Modules, Routines, Types, Templates, Files, Macros, Namespaces, Comments/Pragmas, Statements (C/C++ only)
- http://www.cs.uoregon.edu/research/pdt/





Binary Instrumentation Static binary rewrite Instrumentation code is inserted into the binary before it starts to execute Creates modified executable Dynamic binary instrumentation On-the-fly: Insert, remove, and change instrumentation in the application program while it is running Most flexible (but most complex) technique Parallel programs Coordinated instrumentation of all images needed



Dyninst



- Dyninst is a C++ library for machine-independent
 - $\hfill\Box$ process control and manipulation
 - □ runtime code generation
 - □ and binary patching
- University of Wisconsin and University of Maryland
- Basis for Paradyn, DPCL, and OpenSpeedShop
- Open source
- Supports
 - □ Power/PowerPC (Linux) □ X86 (Linux, BSD, Windows)
 - □ BlueGene/P □ X86_64 (Linux, BSD, Windows)
- http://www.dyninst.org

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Comparison of Techniques (I)



- Source code instrumentation
 - Portable
 - ink back to source code easy
 - Only way to capture "high-level" user abstractions
 - Recompilation necessary for (change in) instrumentation
 - Requires source code to be available
 - (a) Hard to use for mixed-language applications
 - Source-to-source translation tool hard to implement for C++ and Fortran90



Comparison of Techniques (II)



- Binary code instrumentation
 - © / ® The other way around compared to source instrumentation
- Pre-instrumented library / runtime
 - © Easy to use: only re-linking necessary
 - © Can only record information about library / runtime entities
- No single technique is sufficient!
- Typically, combinations of techniques needed!

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CONTENT



- Metrics
- Instrumentation techniques
 - Source code instrumentation
 - · Binary instrumentation
- Instrumentation of parallel programs
 - MPI
 - OpenMP
- Measurement techniques
 - Profiling
 - Tracing

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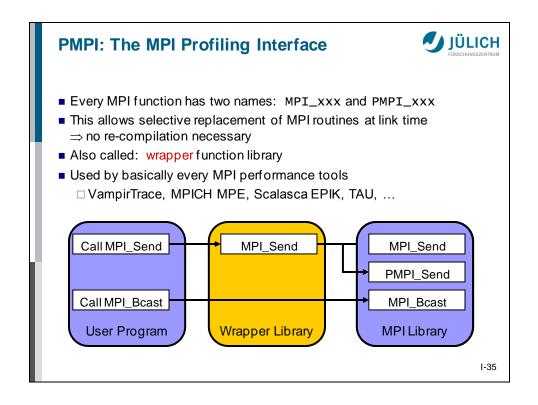
Instrumentation of Parallel Programs	JÜLICH
 User-level constructs Modules / components / Program phases Functions Loops 	
 Constructs of the parallel programming models Message passing MPI, PVM, Threading and synchronization OpenMP, POSIX, Win32, or Java threads, 	
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Instrumentation of User Functions



Ideally: instrumentation by compiler or tool	
☐ Hidden, unsupported compiler options	
(GNU, Intel, IBM, NEC, Sun Fortran, PGI, Hitachi, ???) ☐ TAU Source Code Instrumentor	
☐ TAU Binary Instrumentor (Dyninst)	
□ TAU Virtual Machine Instrumentor (Java, Python)	
 Always works: manually Instrumentation APIs of tools: Scalasca, Vampirtrace, TAU, Scalasca's POMP Directives More details later 	
■ Main problem: selection of relevant constructs	
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PMPI Wrapper Development

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- MPI has many functions! [MPI-1: 130 MPI-2: 320]
 - ⇒ use wrapper generator (e.g., from MPICH MPE)
 - ⇒ needed for Fortran and C/C++
- Message analysis / recording
 - □ Location recording ⇒ use ranks in MPI_COMM_WORLD?
 - □ Data volume

 ⇒ #elements * sizeof(type)
 - □ No message ID □ need complete recording of traffic
 - □ Wildcard source and tag ⇒ record real values
 - □ Collective communication ⇒ communicator tracking
 - □ Non-blocking, persistent communication ⇒ track requests

 - □ One-sided communication?

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OpenMP Monitoring?

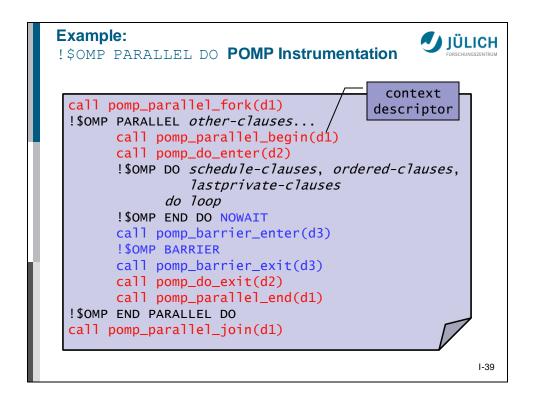


- Problem:
 - □ OpenMP does not define standard monitoring interface
 - □ OpenMP is defined mainly by directives/pragmas
- Solution:
 - □ POMP: OpenMP Monitoring Interface
 - ☐ Joint Development
 - Forschungszentrum Jülich
 - University of Oregon
 - ☐ Presented at EWOMP'01, LACSI'01 and SC'01



"The Journal of Supercomputing", 23, Aug. 2002.





POMP-like Hooks in Production Compilers POMP was the base for the OpenMP instrumentation hooks provided in production compilers Cray Compiling Environment PGI IBM XL compilers These instrumentation hooks are used for performance analysis of OpenMP in production tools CrayPat PGProf Also: New OpenMP ARB sanctioned low-level tool interface http://www.compunity.org/futures/omp-api.html Proof-of-concept implementations by Sun and Intel compilers



POMP Instrumentation Tool



- OpenMP Pragma And Region Instrumentor
- Source-to-source translator to insert POMP calls around OpenMP constructs and API functions



- Implemented in C++
- Supports:
 - □ Fortran77 und Fortran90, OpenMP 2.0
 - □ C und C++, OpenMP 1.0
 - ☐ Additional POMP directives for control and region definition
 - ☐ Used by Scalasca, VampirTrace, TAU, and ompP
 - ☐ Preserves source code information (#line line file)
- Does not support: Instrumentation of user functions

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Current Major OPARI Limitations



- Does not yet support
 - □ Varying number of threads in different parallel regions
 - □ Nested parallelism
 - Latest OpenMP 3.0 standard features like tasking Fixed in OPARI2
- Executed before compiler preprocessor
 - ⇒ issues with macros, conditional compilation, includes!
- Needs special care if building ...
 - ... more than one application in one directory
 - ... applications spread over multiple directories. Fixed in OPARI2

OPARI2: will be available end of 2011



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- Measurement techniques
 - Profiling
 - Tracing

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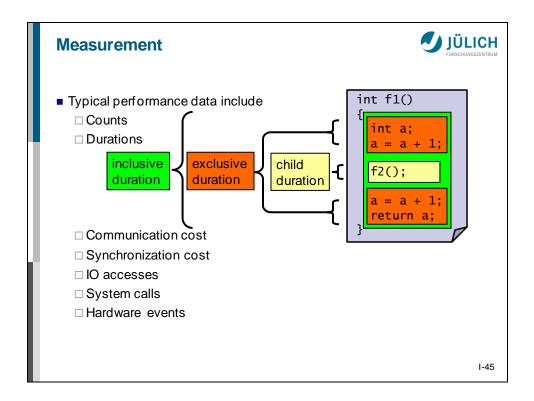
Performance Measurement



- Two dimensions
 - □ When performance measurement is triggered
 - External agent (asynchronous)
 - Sampling
 - Timer interrupt
 - Hardware counters overflow
 - Can measure unmodified executables, very low overhead
 - Internal agent (synchronous)
 - Code instrumentation:
 - Automatic or manual instrumentation
 - ☐ How performance data is recorded
 - Profile ::= Summation of events over time
 - □ run time summarization (functions, call sites, loops, ...)
 - Trace file ::= Sequence of events over time

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Critical Issues Accuracy Perturbation Measurement alters program behavior E.g., memory access pattern Intrusion overhead Measurement itself needs time and thus lowers performance Accuracy of timers, counters Granularity How many measurements How much information / work during each measurement Tradeoff Accuracy ⇔ expressiveness of data



Measurement Methods: Profiling	JÜLICH FORSCHUNGSZENTRUM
 Recording of aggregated information Time Counts Calls Hardware counters about program and system entities 	
□ Functions, call sites, loops, basic blocks, □ Processes, threads	
 ■ Methods to create a profile □ PC sampling (statistical approach) □ Interval timer / direct measurement (deterministic approach) 	
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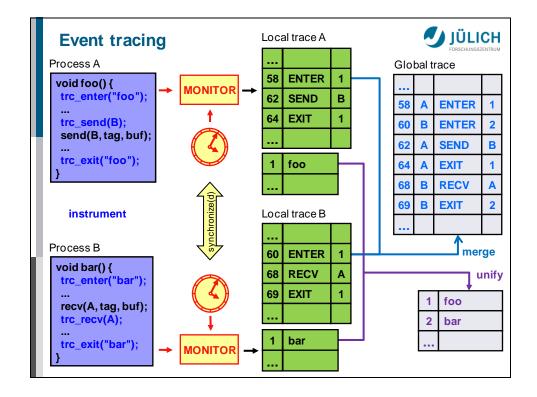
Profiling (2)



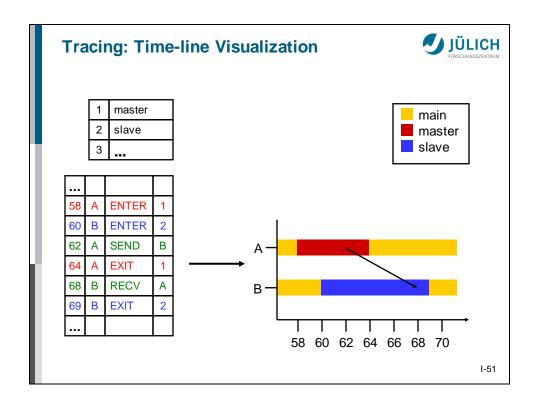
- Sampling
 - ☐ General statistical measurement technique based on the assumption that a subset of a population being examined is representative for the whole population
 - □ Running program is interrupted periodically
 - Operating system signal or Hardware counter overflow
 - Interrupt service routine examines return-address stack to find address of instruction being executed when interrupt occurred
 - □ Using symbol-table information this address is mapped onto specific subroutine
 - □ Requires long-running programs
- Interval timing
 - ☐ Time measurement at the beginning and at the end of a code region
 - □ Requires instrumentation + high-resolution / low-overhead clock



Measurement Methods: Tracing ■ Recording information about significant points (events) during execution of the program □ Enter/leave a code region (function, loop, ...) □ Send/receive a message ... ■ Save information in event record □ Timestamp, location ID, event type □ plus event specific information ■ Event trace := stream of event records sorted by time ■ Can be used to reconstruct the dynamic behavior □ Abstract execution model on level of defined events







JÜLICH Tracing vs. Profiling ■ Tracing Advantages □ Event traces preserve the temporal and spatial relationships among individual events (⇒ context!) ☐ Allows reconstruction of dynamic behavior of application on any required abstraction level Automatic analysis Visualization ☐ Most general measurement technique Profile data can be constructed from event traces Disadvantages □ Traces can become very large ☐ Writing events to a file at runtime can cause perturbation ☐ Writing tracing software is complicated Event buffering, clock synchronization, ... I-52



Trace File Formats



- Current Vampir trace formats
 - □ VTF: family of historical ASCII and binary formats
 - http://www.cs.uoregon.edu/research/paracomp/tau/vtf3-1.43.tar.gz
 - □ OTF: new Open Trace Format
 - http://www.tu-dresden.de/zih/otf/
- TAU performance analysis toolset
 - □ http://tau.uoregon.edu/docs.php#api
- EPILOG: Jülich open-source trace format
 - □ http://www.scalasca.org
- MPICH Multi-Processing Environment (ALOG, CLOG, SLOG, SLOG-2)
 - □ http://www-unix.mcs.anl.gov/perfvis/software/log_format/
- Paraver trace analyzer (BSC, CEPBA)
 - □ http://www.bsc.es/paraver

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No Single Solution is Sufficient! Combination of methods, techniques and tools needed Instrumentation Source code / binary, static / dynamic, manual / automatic Measurement Internal / external trigger, profiling / tracing Analysis Statistics, Visualization, Automatic, Data mining, ...

